Constructions of Large Caps and Progression-Free Sets

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Joint Work with Christian Elsholtz and Benjamin Klahn

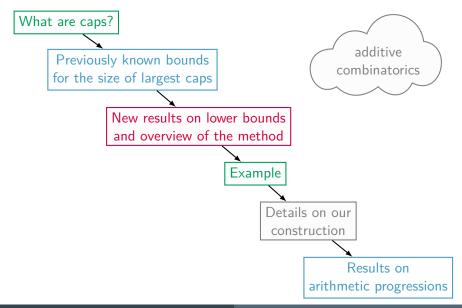
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Plan for the Following 25 Minutes







Definition

An affine (resp. projective) **cap** is a subset of the affine (resp. projective) space in which **no three points lie on a line**.

We mainly consider affine caps in $\mathbb{F}_p^n = (\mathbb{Z}/p\mathbb{Z})^n$ for primes p, and we set

$$C(\mathbb{F}_p^n) \coloneqq \max\{|S| \colon S \text{ is a cap in } \mathbb{F}_p^n\}.$$

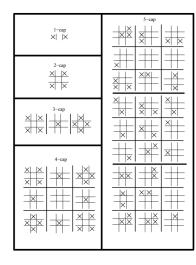
Aim:

construction of large caps in \mathbb{F}_p^n for primes p and arbitrary dimension n

 \hookrightarrow good lower bounds for $C(\mathbb{F}_p^n)$

Since every subset of an affine space can be embedded into the projective space, our lower bounds also hold in the projective case.





- Situation gets complicated very fast.
- It is difficult to find maximal caps in high dimensions.

 \rightsquigarrow bounds

Upper Bounds



For $p \in \{3, 4, 5\}$, we have

"no three points on a line" \iff "no three points in AP".

Theorem

- Ellenberg–Gijswijt (2016): $C(\mathbb{F}_3^n) \leq 2.756^n$,
- Croot-Lev-Pach (2016): $C(\mathbb{Z}_4^n) \le 3.611^n$.

Theorem (Blasiak–Church–Cohn et al. 2017)

We have

 $C(\mathbb{F}_p^n) \leq (J(p)p)^n,$

where

$$J(p) = \frac{1}{p} \min_{0 < t < 1} \frac{1 - t^p}{(1 - t)t^{(p-1)/3}}.$$

Previously Known Lower Bounds

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Best known general constructions so far are "local": take the tensor product of a large cap in small dimension

For a fixed prime p, we have:

Theorem (Bose 1947)

$$\mathcal{C}(\mathbb{F}_p^3)= p^2$$
 and so $\mathcal{C}(\mathbb{F}_p^n)\gg p^{2n/3}$.

Theorem (Edel–Bierbrauer 2004)

$$C(\mathbb{F}_p^6)\geq p^4+p^2-1$$
 and so $C(\mathbb{F}_p^n)\gg (p^4+p^2-1)^{n/6}$.

Theorem (Elsholtz–Pach 2020)

$$C(\mathbb{Z}_4^n) \gg \frac{3^n}{\sqrt{n}}$$
 and $C(\mathbb{F}_5^n) \gg \frac{3^n}{\sqrt{n}}$.

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Theorem (Elsholtz–L 2020+)

$$C(\mathbb{F}_{11}^n) \gg \frac{5^n}{n^{1.5}}, \quad C(\mathbb{F}_{17}^n) \gg \frac{7^n}{n^{2.5}}, \quad C(\mathbb{F}_{23}^n) \gg \frac{9^n}{n^{3.5}},$$
$$C(\mathbb{F}_{29}^n) \gg \frac{10^n}{n^4}, \quad C(\mathbb{F}_{41}^n) \gg \frac{12^n}{n^5}.$$

- exponential improvements for all primes $p \le 41$ with $p \equiv 5 \mod 6$
- "global" and "digit-based" construction based on the method of Elsholtz and Pach for progression-free sets
- basic idea of the construction:

For vectors in the cap,

select a "good" set of digits $D \subseteq \mathbb{F}_p$

and only use these digits for the vectors.

 \hookrightarrow caps of size $(|D| - o(1))^n$

Comparison of the Lower Bounds

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In order to get rid of the dimension in $C(\mathbb{F}_p^n)$, we define

$$c(p) \coloneqq \lim_{n \to \infty} (C(\mathbb{F}_p^n))^{1/n}$$

It is known that the limit exists and $c(p) \in [2, p)$.

p	p ^{2/3}	$(p^4 + p^2 - 1)^{1/6}$	new	improvement
5	2.92401	2.94243	3	1.9562%
7	3.65930	3.67139	3	
11	4.94608	4.95282	5	0.9526%
13	5.52877	5.53418	4	
17	6.61148	6.61528	7	5.8156%
19	7.12036	7.12364	6	
23	8.08757	8.09012	9	11.2468%
29	9.43913	9.44099	\geq 10	\geq 5.9210%
31	9.86827	9.86998	≥ 8	
37	11.10370	11.10505	≥ 10	
41	11.89020	11.89138	≥ 12	\geq 0.9134%

Overview of the Construction

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For a fixed prime p and some set of digits $D \subseteq \mathbb{F}_p$, we consider the set

$$S(D,n) \coloneqq \left\{ (a_1, \ldots, a_n) \in D^n \, \middle| \, \forall d \in D \colon a_i = d \text{ for } \frac{n}{|D|} \text{ values of } i
ight\}.$$

We call D good if S(D, n) is a cap for all appropriate $n \in \mathbb{N}$. By Stirling's formula, we obtain

$$|S(D,n)| = \prod_{\ell=0}^{|D|-1} \binom{n-\frac{\ell n}{|D|}}{\frac{n}{|D|}} \sim \frac{c|D|^n}{n^{\delta}}$$

with

$$\delta = rac{|D|-1}{2}$$
 and $c = rac{1}{\sqrt{1-\delta/|D|}} \Big(rac{|D|}{2\pi}\Big)^{\delta/2}$

.

Connection and Difference to APs



Three-term arithmetic progressions are solutions of the equation

$$x - 2y + z = 0. \tag{(\star)}$$

Three points x, y, $z \in \mathbb{F}_p^n$ are **not collinear** if and only if $ax + by + cz \neq 0$ for all $(a, b, c) \in \mathbb{F}_p^3 \setminus \{(0, 0, 0)\}$ with a + b + c = 0.

Without loss of generality, we can assume a = 1 and $b \notin \{-1, 0\}$.

Three points x, y, $z \in \mathbb{F}_p^n$ are **not collinear** if and only if

$$x + by + (-b - 1)z \neq 0$$
 for all $b \in \mathbb{F}_p \setminus \{-1, 0\}$. $(\star\star)$

 \hookrightarrow still p-2 equations to consider

Idea: Apply the method of Elsholtz and Pach not only to (\star) , but also to the other equations $(\star\star)$ corresponding to "weighted progressions".

~ much more involved

Example: p = 11



We choose

- the digit set $D = \{0, 1, 3, 4, 5\}.$
- If D is good, then this implies

$$C(\mathbb{F}_{11}^n)\gg \frac{5^n}{n^2}.$$

Equivalent equations:

•
$$\{x - 2y + z = 0, x - 10y + 9z = 0, x - 6y + 5z = 0\},$$

• $\{x - 3y + 2z = 0, x - 7y + 6z = 0, x - 9y + 8z = 0, x - 5y + 4z = 0, x - 8y + 7z = 0, x - 4y + 3z = 0\}.$
• $x - 2y + z = 0:$
 $P_{-2}(D) = \{(1, 3, 5), (3, 4, 5), (5, 3, 1), (5, 4, 3)\}$
 $\hookrightarrow \{(3, 4, 5), (5, 4, 3)\} \rightarrow \emptyset$
• $x - 3y + 2z = 0:$
 $P_{-3}(D) = \{(1, 0, 5), (1, 3, 4), (1, 4, 0), (3, 0, 4), (3, 1, 0), (4, 1, 5), (4, 5, 0), (5, 0, 3)\} \rightarrow \emptyset$

Finding Good Digit Sets (I)



We fix $b \in \mathbb{F}_p \setminus \{-1, 0\}$ and $D \subseteq \mathbb{F}_p$, and set

$$P_b(D) = \left\{ (x,y,z) \in D^3 \, \Big| \, x + by + (-b-1)z = 0 \right\} \setminus \left\langle (1,1,1) \right\rangle.$$

Assume that there is some $n \in \mathbb{N}$ with $|D| \mid n$ such that there are 3 points

$$x = (x_1, \ldots, x_n)^{\top}, \ y = (y_1, \ldots, y_n)^{\top}, \ z = (z_1, \ldots, z_n)^{\top} \in S(D, n)$$

which satisfy x + by + (-b - 1)z = 0.

→ **introduce variable** χ_v for each $v = (v_1, v_2, v_3) \in P_b(D)$ which describes the number of occurrences of v in the components of x, y, z, i.e.,

$$\chi_{v} = |\{i \in \{1, \ldots, n\} | (x_{i}, y_{i}, z_{i}) = v\}|.$$

Since every digit d in D has to occur the same number of times, we find

$$\sum_{\substack{\nu \in P_b(D) \\ \nu_1 = d}} \chi_{\nu} = \sum_{\substack{\nu \in P_b(D) \\ \nu_2 = d}} \chi_{\nu} \text{ and } \sum_{\substack{\nu \in P_b(D) \\ \nu_1 = d}} \chi_{\nu} = \sum_{\substack{\nu \in P_b(D) \\ \nu_3 = d}} \chi_{\nu}.$$

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Finding Good Digit Sets (II)



$$\sum_{\substack{\nu \in P_b(D)\\\nu_1 = d}} \chi_{\nu} = \sum_{\substack{\nu \in P_b(D)\\\nu_2 = d}} \chi_{\nu} \quad \text{and} \quad \sum_{\substack{\nu \in P_b(D)\\\nu_1 = d}} \chi_{\nu} = \sum_{\substack{\nu \in P_b(D)\\\nu_3 = d}} \chi_{\nu} \quad (\star)$$

S(D, n) does not contain x, y, z with x + by + (-b-1)z = 0 for \iff System (*) has no non-trivial non-negative integral solution $\chi = (\chi_v \mid v \in P_b(D)).$

Hence, to show the "goodness" of some D, one has to ensure that

 $\mathcal{P} = \{ \chi \in \mathbb{Z}^{\ell}_{\geq \mathbf{0}} \, | \, \mathbf{A} \cdot \chi = \mathbf{0} \}$

is empty, where the matrix A represents (\star) .

→ integer programming

- Appropriate software is available. 😊
- Checking the emptiness of $\mathcal P$ is NP-complete. $\ensuremath{\mathfrak{O}}$

\rightsquigarrow simpler conditions required

Arithmetic Progressions



Let $r_k(\mathbb{F}_p^n)$ denote the size of the largest progression-free set in \mathbb{F}_p^n .

Theorem (Lin–Wolf 2010)

If $k \leq p$, then we have

$$r_k(\mathbb{F}_p^n) \ge (p^{2(k-1)} + p^{k-1} - 1)^{\frac{n}{2k}} \approx p^{\frac{(k-1)n}{k}}.$$

Theorem (Elsholtz–Pach 2020)

For $p \ge 5$ and some explicitly given constant d_p , we have

$$r_3(\mathbb{F}_p^n) \geq \frac{d_p}{\sqrt{n}} \Big(\frac{p+1}{2}\Big)^n.$$

Theorem (Elsholtz–Klahn–L 2020+)

 $r_5(\mathbb{F}_{23}^n) \gg (17 - o(1))^n$ $r_7(\mathbb{F}_{29}^n) \gg (24 - o(1))^n$ (improving on 12.28ⁿ) (improving on 17.92ⁿ)

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Thank you for your attention!